Use of analogy in a production system architecture

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troduction

analogical problem solving was fundamental in the skill acquisition developed in response to these ruminations. architecture. This is the first report of the new system that we have how analogy takes place and how it fits into a production system no explanation as to why students were using analogy as their primary tecture. It was also ad hoc in the sense that the architecture offered ever, the process of simulating this was awkward within that archi-Anderson, Pirolli, & Farrell, 1988; Pirolli & Anderson, 1985). Howwithin the ACT* architecture (Anderson, Farrell, & Sauers, 1984; problems. We did a number of simulations of this problem solving books or to examples from more familiar domains for solving these solving. Students were always resorting to examples from their textdomains we were studying: programming and mathematical problem system architecture. A couple of years ago we became persuaded that This is a report on our development of analogy within a production problem-solving method. This led us down the path of thinking about

The term analogy is used in multiple senses. We are concerned with how analogy is involved in problem solving and in skill acquisition. That is, how do people call up an analogous experience to help them solve new problems? This experience can come from their own past (when analogy is sometimes called repetition), it might come from looking at the behavior of another (in which case it is sometimes called imitation), or it might come from adapting an example given in a textbook or some other expository medium (in which case it is sometimes called copying). The source for the analogy can be either an explicit experience or a generic or schemalike representation. It also might be from the "same" domain as the problem or from a different domain. It may involve generating a solution to a problem or understanding a solution. This range of phenomena has not always been

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organized as a single psychological kind. However, they flow from a single mechanism in our theory.

Our mechanism does not address the more literary or mnemonic uses of analogy. For instance, we have nothing to say about the typical use of the solar system to understand the atom. The typical uses of that analogy do not involve any problem solving in our analysis. This is not to say that these are not interesting psychological issues; it is just to define the domain of reference for our chapter.

Our analysis of problem solving sees knowledge organized into function and form. Generating a solution involves producing a form that will achieve a desired function. Understanding a problem solution involves assigning a function to an observed form. Analogy is a mechanism for achieving the function-form relationships in domains where one has not yet acquired skills for doing so. Skill acquisition involves acquiring productions that circumvent the need to go through the analogy mechanism.

In this chapter we will first describe our knowledge representation, which is factored into a function—form structure. Next we will describe how the analogy process operates on this to fill in missing function and form. Then we will discuss the knowledge compilation process, which operates on the trace of an analogy process to produce new productions (i.e., acquire skill). Finally, we will discuss some issues associated with placing this process into a more general framework of skill acquisition. Before doing any of this, however, we would like to explain a little about the PUPS production system in which the analogy mechanism is implemented.

The PUPS architecture

PUPS (for PenUltimate Production System) is intended as an implementation of the theoretical successor to the ACT* theory (Anderson, 1983). This new theory does not yet really exist. We are working on this production system as a means of developing that theory. Therefore, the extent of the difference between the PUPS-based theory and its ACT* predecessor is not clear. Nevertheless, there are a series of problems with the ACT* theory that PUPS intended to remedy.

1. The flow of control in ACT* was implemented by means of a goal stack that yielded top-down, fixed-order traversal of goals. There were two basic problems with this: (a) People can be more flexible in their flow of control. For instance, as Van Lehn (1983) has noted, children doing multicolumn subtraction problems will shift opportunistically among the preferred right-to-left processing of columns,

Such flexibility is the rule rather than the exception when we watch geometry students develop proofs. (b) The ACT* theory leads to unreasonable expectations about people's ability to remember goals. It also does not predict the fact that people do much better at remembering goals when there is a concrete residue of their problem solving such as marks on a page than when they must solve these problems is to abandon the use of goal stacks as separate control structures and let the problem solution itself hold the goals. This ties memory for solution to memory for pending goals as it appears to be. It also makes goals declarative structures and so permits greater flexibility in selection of the goal to follow.

2. In ACT*, generalization and discrimination were automatic processes calculated by the architecture on productions. As a consequence, they were not open to inspection or conscious control. In contrast, it is argued (Anderson, 1986) that people have some access and control over their induction. As we shall see when we discuss analogy and skill acquisition, the new architecture of PUPS enables this less automatic sort of skill induction.

3. Analogy itself was cumbersome to calculate in ACT* because there were not the right representational or processing primitives. In contrast, in every domain of skill acquisition we studied, we found that analogy was the prime mechanism by which students learned to solve new problems. Therefore, it seemed that the architecture had to be reconfigured to permit more natural computation of analogy.

PUPS knowledge representation

Our theory is built around certain assumptions about the organization of knowledge. The basic assumption is that knowledge is represented in schemalike structures. The three obligatory slots for such a structure are a category (isa) slot that specifies the category of which the structure is an instance, a function slot that specifies what function the structure fulfills, and a form slot that specifies the form of the structure. Below is our representation of the LISP function call (+ 2 3):

structure1
isa: function-call
function: (add 2 3)
form: (list + 2 3)
context: LISP
medium: CRT-screen

precondition: context: LISP

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This asserts that this is a function call, its function is to add 2 and 3, and its form is literally a list of a +, followed by a 2, followed by a 3. Note there are three optional slots for this structure. The first asserts that this function call is being used in the context of LISP, the second asserts that this is being entered on a CRT screen, and the third asserts that it is critical that the context be LISP for the form to achieve the specified function. Entering this form in other programming languages would not have this effect. On the other hand, the medium is an accidental property of this example and so is not listed as a prerequisite.

The function, form, and prerequisite slots have a certain implicational semantics, which it is useful to set forth at the outset. This implicational semantics is that the form and the prerequisites imply the function.

We can represent this as

(list + 2 3) & context (LISP) \rightarrow (add 2 3)

As we shall see, analogy creates function and form slots that plausibly create such implications.

There are two complications that greatly enhance the expressive power of our knowledge representation. One is that a structure can serve multiple functions and so have multiple values in its function slot. The second is that one structure can be embedded in another. Consider the following structures that represent the LISP code:

(cons (cadr lis) (cons (car lis) (cddr lis)))
structure2
isa: function-call

function: (reverse elem1 elem2)
(insert elem2 lis2)
(find lis1)

form: (list cons structure3 structure4)

structure3
isa: function-call
function: (extract-second lis)
(find elem2)
form: (list cadr lis)

structure4
isa: function-call
function: (insert elem1 list3)
(find lis2)
form: (list cons structure5 structure6)

```
elem2
                                                                                                                                                                                                                                                                                                                                                                                                                                                                           elem I
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                               structure6
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                   structure5
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         function: (extract-second-tail lis)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                           isa: function-call
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         torm: (list cadr lis)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                             function: (extract-first lis)
                                                                                                                                                                                                        tunction: (second-tail lis)
                                                                                                                                                                                                                                                                                                              function: (second lis)
                                                                                                                                                                                                                                                                                                                                 isa: element
                                                                                                                                                                                                                                                                                                                                                                                                                                                        isa: element
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                     form: (list cddr lis)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                  isa: function-call
                                                                 function: (value-of structure4) (first-tail lis1)
                                                                                                        isa: list
                                                                                                                                                                                                                              isa: list
                                                                                                                                                                                                                                                                                                                                                                                                                                     function: (first lis)
function: (value-of structure2)
                     isa: list
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                        (find lis3)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                               (find elem1)
                                                                                                                                                                                                                                                                                                                                                                             (value-of structure5)
                                                                                                                                                                                                                                                                                                                                                                                                  (second lis1)
                                                                                                                                                                                                                                                                                                                                                                                                                     (first lis2)
                                                                                                                                                                                         (second-tail lis1)
                                                                                                                                                                                                                                                                          (value-of structure3)
                                                                                                                                                     (value-of structure6)
                                                                                                                                                                        (first-tail lis2)
                                                                                                                                                                                                                                                                                               (hrst list)
```

This is a fairly elaborate representation of this little piece of LISP code, and we do not mean to imply that all LISP programmers always have as rich a representation. However, we will see that success in problem solving can turn on the richness and correctness of such structure representations.

Analogy

There are number of distinct cases of analogy that need to be considered. The one that is most basic in our work is the filling in of a

example, consider the structures below: missing form slot in a structure that has a filled-in function slot. For

structurex function: (extract-first (a b c)) form: ??? isa: function-call

structurey isa: function-call form (list car '(p q r s)) function: (extract-first (p q r s))

achieves an analogous goal. but missing form. Structurey is a structure with a form filled in that how we represent a goal in PUPS - a structure with a filled-in function element of the list (a b c), but there is no form specification. This is Structurex above has a functional specification that it gets the first

The analogy we solve to fill in the form slot is the following:

function(structurey) : form(structurey) :: function(structurex) : ???

or.

(extract-first(pqrs)):(listcar'(pqrs))::(extract-first(abc)): ???

tween elements of the form slot of the two structures: structurex. This is produced by creating the following mapping be-The solution of course is to put (hst car '(a b c)) in the form slot of

list car \rightarrow (a b c) → car **→ list**

and which must be mapped onto new symbols from the target domain. explain how we decide which symbols can be mapped to themselves special and maps onto itself. The symbols car and 'are mapped onto tures. The first element in the form slot (list, in this case) is always themselves and $(p \ q \ r \ s)$ is mapped onto $(a \ b \ c)$. Momentarily, we will correspondence established between the functions of the two struc-Three of these elements map onto themselves, and the other uses a

are available for the mapping in the analogy. are put into correspondence (just $(p \ q \ r \ s)$ and $(a \ b \ c)$ in this case) and (extract-first, in this case) must correspond. The remaining elements to be considered successful the first elements in the two function slots The mappings are obtained from the function slots. For an analogy

> structurey as: tionship between function and form we can represent the content of in this analogy. Going back to our implicational analysis of the rela-It is worthwhile to identify the inductive inference that is contained

(list car '(p q r s)) \rightarrow (extract-first (p q r s))

generalization: What we have done in making the analogy is to make the

(list car 'x) \rightarrow (extract-first x)

called the no-function-in-content principle. That is, the actual content the structurex as a special case. This inductive inference step we have thus, variabilizing the nonterminal (p q r s). Given this, we can derive function achieved by the form in which they appear. identities of nonterminal symbols like (p q r s) are not critical to the

termines which constants can be left and which must be mapped in are simply mapped onto themselves in the target domain. This demapped in the form slot as they are in this example (car and '). They for all terms in the function slot. It is all right if terms are left unway, the necessary and sufficient condition is that mappings be found tiation of the antecedent we can derive the consequent. Said another bilized (except the first special terms). Then, by producing an instanbuilding up a new form. that all the terms in the consequent of the above implication be varia-Note that the necessary and sufficient condition for an analogy is

of how someone encodes a call to the LISP function quotient by analogy by simply comparing the function slots of the target and the goal to the LISP function difference: However, this is not always possible. Consider the following example In the preceding example the mapping for (pqrs) could be obtained

Example

ısa: tunction-call

function: (perform subtraction 6 2) form: (list func1 6 2)

isa: LISP·function

function: (implement subtraction)

form: (text difference)

isa: function-call function: (perform division 9 3) form: ??

isa: LISP-function function: (implement division) form: (text quotient)

This represents the knowledge state of someone who knows what both difference and quotient do but has seen only the syntax for difference and is trying to figure out the syntax for quotient. In looking at the function slot of this example to make the correspondences for the analogy, we get the following:

```
subtraction \rightarrow division 6 \rightarrow 9 2 \rightarrow 3
```

However, this does not provide a specification about how to map the term func! that appears in the form slot of the example. But if PUPS looks at the function slot of func!, it will find that it implements subtraction. Knowing that subtraction corresponds to division, it knows it wants the name of a function that implements division. Because it knows quotient is such a function, it writes (LIST quotient 9 3) into the form slot of goal.

If we return to our implicational analysis we can see what assumptions are involved in making this step. The analysis of the example would be

```
(list difference y z) \rightarrow (perform subtraction y z)
```

We have replaced the constants 6 and 2 by y and z to reflect the nofunction-in-content principle. However, this will not match the function slot of *goal* because of the unmapped element subtraction. However, if we replace difference by its functional specification we get

```
(LIST (implement subtraction) 6 2) \rightarrow (perform subtraction 6 2)
```

Applying the no-function-in-content principle to this example, we get the following variabilized expression:

```
(list (implement x) y z) \rightarrow (perform x y z)
```

This essentially says, if we create a list of the function name that implements an operation and two arguments, we will apply that operation to the arguments. In replacing difference by its functional specification we are assuming that the critical thing about difference was its functionality and not its identity. If there is another function that implements subtraction (as there is), then it would work the same way. This is another inductive principle, which we have called sufficiency of functional specification.

isa: LISP.function

These two inductive principles, no function in content and sufficiency of functional specification can in combination produce rather elaborate problem solutions by analogy. To take one example, we were able to get the mechanism to code the function summorial by (10 iterations of) analogy to the function factorial, both of which are given below:

```
(defun factorial (n)
  (cond ((zerop n) 1)
   (t (times n (factorial (sub1 n))))))
(defun summorial (i)
  (cond ((zerop i) 0)
   (t (plus i (summorial (sub1 i))))))
```

Factorial calculates the product of the integers to n whereas summorial calculates the sum. Both are given recursive definitions.

Although we haven't space to go through this whole analogy process in detail, it is worth going through a few key steps. First, the following gives some of the original encoding of the example and the problem:

```
sum.goal
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                        tact.algorithm
                                                                                                                                                                                                                              fact.args
                                                                                                                                                                                                                                                                                                                                               tact.name
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                               isa: LISP.function
                                                                                                                                                                                                                                                                                                                     isa: function.name
                                                                                                                                                                                                                                                                                                                                                                                                                                   isa: variable
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                             form: (compute product (range zero fact-arg))
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                      tunction: (algorithm-for fact.ex)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                          ısa: algorithm
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                               form: (LIST defun fact.name fact.args fact.body)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                      function: (calculate fact.algorithm)
                                                                                                                                                                                                   ısa: param.list
                                                                                                                                                                                                                                                                                                                                                                                                          function: (parameter-for fact.algorithm)
                                        function: (body-for fact.ex) form: (LIST cond case 1 case 2)
                                                                                        isa: LISP.code
                                                                                                                                                                              function: (enumerate-args fact.ex)
                                                                                                                                                                                                                                                                                             function: (name-of fact.algorithm)
                                                                                                                                                                                                                                                                                                                                                                                     torm: (text n)
                                                                                                                                                     form: (list fact.arg)
                                                                                                                                                                                                                                                                     torm: (text factorial)
```

```
function: (calculate sum.algorithm)
```

sum.algorithm ısa: algorithm form: (compute sum (range zero sum.arg)) function: (algorithm-for sum.goal)

isa: variable function: (parameter-for sum.algorithm)

form: (text 1)

isa: function.name function: (name-of sum.algorithm)

form: (text summorial)

of the summorial code. To follow the analogy process, it starts with a representation of the implicational structure of fact.ex: The first task that is solved by analogy is to determine the top level

(list defun fact.name fact.args fact.body) → (calculate fact.algorithm)

ficiency-of-functional-specification principle) we get: replacing terms with their functional specification (applying the suf-In this form it does not allow for any variabilization; however, by

(list defun (name-of fact.algorithm) fact.args fact.body) → (calculate fact.algorithm)

which will invoke elaboration of fact.args and fact.body and lead to a for fact.ex). This creates another element fact.ex on the right-hand side we embellish fact.algorithm with its functional specification (algorithm sufficiency-of-functional-specification principle can apply. In this case we have a third principle, which we call maximal functional elaboration and it would change its behavior. Clearly, this is not acceptable. Thus, who thought that it was sufficient to change the name of function directly copied to the target structure. This would represent a subject torm that can have greater variabilization: implication will be embellished with their functions as long as the (left-hand side), arguments in the function (right-hand side) of the This principle is that, if there are unvariabilized terms in the form Note fact.args and fact.body would not be variabilized and would be This expression could be variabilized and mapped to the target

(list defun (name-of fact algorithm) (enumerate-args fact ex) (body-for fact.ex)) (calculate fact.algorithm = (algorithm-for fact.ex))

> ciple to give: This can now be variabilized using the no-function-in-content prin-

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(list defun (name-of x) (enumerate-args y) (body-for y) \rightarrow (calculate x = (algorithm-for y))

Instantiating this with the summorial function we get

(list defun (name-of sum.algorithm) (enumerate-args sum.goal) → (calculate sum.algorithm = (algorithm-for sum.goal) (body-for sum.goal))

to fill the last two slots. So the following structure is inserted into the However, it has to create new PUPS structures, goal.args and goal body, analogy will use the existing structure rather than creating a new one. morial function. In cases where a structure serving the function exists, and the body for the function. There is already a name for the suma list consisting of defun, the name of summorial, the list of arguments, form slot of sum.goal: Thus we need to fill in the form slot of the summorial function with

(list defun sum.name goal.args goal.body),

and goal.body. and the following two data structures are created to describe goal args

goal.args

isa: param.list

function: (enumerate-args sum goal)

form:???

goal.body

isa: LISP. code

function: (body-for sum.goal) form: ???

all the levels of the recursive code for summorial. ogy. In our simulation, these goals were themselves solved by analogy. we have spawned two subgoals as the by-product of solving this anal-The goal.body structure spawned a rich set of goals corresponding to These structures act as goals to invoke further problem solving. Thus

¹ (summorial (sub1 i))). The example and the goal at this point were generate the code that would correspond to the recursive action: (plus ogy problem solving, one of which occurred when it came time to There are a couple of interesting moments in the subsequent anal-

fact.rec.value

isa: LISP.expression

form: (list times fact.arg fact.recursive.call) function: (recursive-value fact.algorithm)

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isa: LISP.primitive

function: (compute product)

form: (text times)

act.recursive.call

isa: recursive.call lunction: (recurse-on fact.ex)

form: (list fact.name fact.rec.arg)

tunction: (recursive-value sum.algorithm)
form: ??? isa: LISP.expression

sum.rec.value

Again, the implicational structure of the example, fact.rec.value, is:

(list times fact.arg fact.recursive.call)

→ (recursive-value fact.algorithm)

The terms in this can be rewritten to become

(list (compute product) (parameter-for fact.algorithm)

(compute product (range zero (parameter-for fact.algorithm)))) \rightarrow (recursive-value fact.algorithm = (algorithm-for fact.ex) &

that corresponds to sum in the specification of summorial. its form slot as well as its function. This gives us the term product What is noteworthy about this is that we have rewritten factorial by

Variabilized, the expression becomes:

(list (compute x) (parameter-for y) (recurse-on z)) \rightarrow (recursive-value y = (algorithm-tor z) &

(compute x (range zero (parameter-for y)))),

which, applied to the summorial case, causes us to create the form:

(list plus sum.arg goal.recursive.call)

ognized as the argument for summorial; and goal.recursive.call is a goal structure created to calculate the form of the recursive call Plus is recognized as the function that computes sum; sum.arg is rec-

Extending the analogy model

Analogical filling of function slots

example to a target using correspondences set up in mapping the in mapping the form. In this we are dealing with a situation where is, to map a function from an example using correspondences set up function. However, analogy can also be used to do the reverse; that So far we have discussed how analogy is used to map a form from an

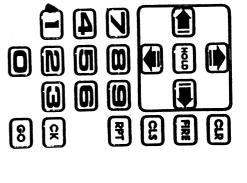


Figure 9.1. Keypad on the toy tank used by Shrager (1985).

what the keys with the down-arrow and left-arrow did. Below we have the up-arrow moved the tank forward, and they had to figure out the keypad in Figure 9.1. They determined that the key labeled with by Shrager (1985), subjects were presented with a toy tank that had the form slot is filled but the function slot is not. In an experiment PUPS structures that purport to represent their states of knowledge

isa: button

form: (labeled up-arrow) function: (move forward)

up-arrow

ısa: symbol

function: (points forward)

form: (image thing1)

problem l

isa: button

form: (labeled down-arrow) function: ???

down-arrow ısa: symbol function: (points backward)

problem2

form: (image thing2)

function: ??? isa: button

form: (labeled left-arrow)

isa: symbol function: (points leftward) form: (image thing3)

The example is encoded as an up-arrow with the further information that an up-arrow is a symbol that conventionally means *forward*. The functions of the other two buttons are not represented, but we have represented the conventional knowledge that down-arrows symbolize backward and left-arrows left.

We can represent the knowledge encoded by the example by the following variabilized expression:

(labeled (points x)) \rightarrow (move x)

This implication represents the operation of the same heuristics as we saw before. This can now be instantiated for one of the problems as:

(labeled (points backward)) \rightarrow (move backward)

Hence, we can infer that the function of the problem1 button is to move backward. Similarly, we can infer that the function of the problem2 button is to move left. As it turned out, only the first inference was correct. The left-arrow button did not actually move the tank in the left direction but only turned it in that direction. This is an example of where the no-function-in-content assumption was volated. Some buttons moved the tank in the specified direction, and some turned. One simply had to learn which did which. The actual identity of the direction determined the function of the button. This just proves that analogy has the danger of any inductive inference. The important observation is that human subjects also made this misanalogy.

Refinement

Though each structure can have one form slot, it can serve multiple functions. This leaves open a third and important kind of analogy: the filling in of a second function by analogy to a first.² This kind of analogy is very important in problem solving because the key to finding the form for solving a problem might be to represent correctly the function of a structure. This process of producing a new functional slot (or new views on the structure) is what we call refinement. Its usefulness is illustrated in the following example, which comes quite close to what subjects originally do when they initially reason about how to call LISP functions.

Consider how someone might analogize from an example of car to determining the value of (cdr'(xy)). Suppose, they had an example of (car'(abc)) = a. They might encode this:

example
isa: LISP.call
function: (show-value expl value!)
form: (string expl = value!)
expl
ica: LISP expression

isa: LISP expression
function: (calculate first arg1)
form: (list car ' arg1)
ar
isa: LISP function
function: (calculate first)
form: (text car)
alue1

The problem would be encoded:

problem
isa: LISP.call
function: (show-value exp2 goal2)
form: (string exp2 = goal2)

exp2
isa: LISP.expression
function: ???
form: (list cdr ' arg2)
cdr
isa: LISP.function
function: (calculate rest)

goal2
isa: sexpression
function: (value-of exp2)
form: ???

form: (text cdr)

arg2 isa: list

slot of goal2 to create a second functional description of goal2: (rest out the function of arg2 or exp2 from the cdr example, and the goal arg2). This is an example of goal refinement; goal2 now has the form: analogy between the two function slots of value 1 and the one function filling in the form slot of goal2 remains unsatisfied. We can form an exp2 is filled in (calculate rest arg2). However, the important task of the value of the cdr example. Using form-to-function analogy, the is to figure out the form of goal2, which corresponds to figuring out function slot of arg2 is filled in as (hold (x y)), and the function slot of This represents an encoding where the subject has not yet figured

function: (value of exp2) isa: sexpression (rest arg2)

the system can look to its other, non-LISP knowledge for an analog. no illustration of the rest relationship. However, given this refinement sending out invitations to a list of people to attend a meeting: So we might have the following encoding of a state when we were This cannot be solved by analogy to the car example, because it has

new3 function: (hold (John Mary Tom))
(invitees meeting1) (hold (Mary Tom)) form: (list Mary Tom) tunction: (rest arg3) form: (list John Mary Tom) isa: list

By analogy to this example, we can fill in the form slot of goal2 by

tion. Now we can call on an extradomain analogy to solve the problem. cannot solve but refine its function so that it has a different descripthan if a single model has been mapped. to solve a problem. The resulting solution is inherently more "novel" This is one example of many where multiple models can be combined be called creative problem solving. We start with a problem that we This example illustrates the characteristic scenario of what might

Selecting examples for use in analogy: spreading activation

ing" will consider discrimination of overly general knowledge. ogy with learned productions, and the section "Discrimination Learnselected for use as models in the analogy, the section "Knowledge sections of this chapter. This section will consider how examples are complications to this, which we will discuss in the next three major novel situation can solve a problem by analogy. There are three major Compilation" will consider how knowledge compilation replaces anal-What we have discussed so far is a method by which a person in a

trolled in the PUPS system by productions of the form: in the the process of matching productions. Basically, analogy is conposal put forth in Anderson (1983, chap. 5), that analogs are chosen Our idea of how analogs are selected is no different than the pro-

F serves that function and there is a model structure containing a form that specified function there is a target structure needing a form serving a

try to map the model form to the target forn

same function. It is possible to have more heuristic versions of this of a production. This is a "bare-bones" selection production in that certain cases the set of possible candidates might be reduced. choosing from multiple possible candidates. It means only that in the only criterion for selecting a model is that the model serve the However, this complication does not really eliminate the problem of production, which used domain-specific tests to look for likely analogs. Thus analogy is an action that can be called on the right-hand side

senting this production is: of the above production is matched. The actual PUPS code repre-The critical issue for selecting a model is how the second condition

(p draw-analogy analogy! = model = target form) = target: isa = object = model: isa = object function (=rel) form nil form (group! <> nil) function (=rel)

for any structure (which will be bound to = model) for which the first The terms = model and = rel are variables. This production looks

term in the function slot is the same as for the target goal (this test is enforced by the appearance of the same variable in the function slots of the goal and the model). Of course, the arguments of the function slot can be different and will be put into correspondence for purpose of analogy. What is important is that the first, predicate, terms are the same, indicating similar function. PUPS will calculate all instantiations of this production and so find all the possible models. The issue is how it selects among these instantiations. This is the issue of conflict resolution.

in the ACT* model. Thus, basically, the most active model structure that the subject is likely to use recent and frequently studied examples verging on it from active elements. The first factor, strength, means strength of that structure and the number of network pathways conabout this suggestion. Activation of a structure basically reflects the will be the one selected. There are a lot of things that seem right tivation-based pattern-matching mechanisms that solved this problem among possible models is quite superficial the functionality of the problem features, the criterion for selecting contrast to the analogy computation itself, which carefully examines geometry but without the benefit of careful data like Ross's. Thus, in We have argued a similar point (Anderson, 1983) for the domain of example, whether the features are relevant to the problem or not. with the current problem determined the probability of selecting the lection. He found that the number of features an example shared volume) points to the importance of feature overlap in analogy semains like LISP and geometry. Research such as Ross's (1984, this pretty apparent in our research on analogy in problem-solving do-The dominance of recency and frequency in example selection is Our view is that conflict resolution is determined by the same ac

Knowledge compilation

Knowledge compilation is motivated to deal with the computational costs of problem solving by analogy. Analogy is expensive to compute and requires having an example at hand and holding a representation of the example in working memory. Knowledge compilation is the process that inspects the trace of analogy and builds productions that directly produce the effect of analogy without having to make reference to the example. Our view is that this knowledge compilation process occurs simultaneously with the first successful analogy. Subsequent occasions where the knowledge is required show the benefit of the compiled production. This corresponds to the marked im-

provement in speed and accuracy from first trial to second in a typical problem-solving situation. We typically see more than 50% improvement and a concomitant marked decrease in any verbalization, indicating that the analogy is being computed (Anderson, 1982, 1987b). Thus analogy is something done only the first time the knowledge is needed (see also Holyoak & Thagard, this volume).

We have adapted the knowledge compilation process (Anderson, 1986) to operate in PUPS. It compiles productions from the trace of the analogy process. Thus, after writing the summorial function by analogy to factorial, PUPS compiled a set of productions that represented the transformations being computed by analogy. For instance, it formed the following production:

the goal is to write a LISP function y
which calculates an algorithm x
create the form (list defun name args body)
where name is the name of algorithm x
and args enumerate the arguments for function y
and body is the body for function y

This is basically an embodiment of the abstract implicational structure that we extracted in doing the analogy:

(list defun (name-of x) (enumerate-args y) (body-for y)) \rightarrow (calculate x = (algorithm-for y))

Thus, knowledge compilation stores away the implication that we had to induce to perform the analogy. The availability of that implication saves us from having to calculate it a second time.

The other thing that knowledge compilation will do is to collapse a number of steps of problem solution into a production that produces the same effect in a single step. To consider an example, suppose we start with the goal to code the second element of a list:

joal isa: function-call function: (extract-second lis) form: ???

Suppose this has its function slot refined by analogy or an existing production to have the additional specification (extract-first lis1) where lis1 is the tail of the list. The structure lis1 is specified as:

isa: list function: (extract-tail lis) form: ???

the function slot of goal and filling in the two structure slots: for list becomes (list cdr lis). In all, three steps were involved, refining existing productions. The form slot for goal becomes (list car list) and The form slots for goal and list can then be solved by analogy or

A production can be composed to summarize this computation:

F THEN which calculates the second of lis y the goal is to code a function call x

create a form (list car struct) where struct is a function call that calculates the tail of

and has form (list cdr y)

purposes of learning. inspects declarative traces and not procedural traces. Declarative Procedural traces were an invention in earlier theory just for the PUPS structures are already there and are by definition inspectable It is more realistic to propose that the learning mechanism simply ACT* theory. It can be extracted directly from the PUPS structures. we do not need a procedural trace to define compilation as in the volved without actually inspecting the productions that fired. That is, pilation of productions can be defined on the PUPS structures in-PUPS structures goal and lis1. This is an interesting observation: Com-This is just the composition of the implication structures of the two

compositions first noted by Lewis (1981). modified composition to deal with the problem of noncontiguous contiguous fragment of the problem solution. In this way we have be invoked by a single ACT* goal. They only have to produce a the productions to be composed fire contiguously in time or that they Another advantage of this is that we do not have to require that

Discrimination learning

and so we do not have the same inflexible mechanism that existed in the PUPS encodings, one can change the direction of generalization, out rationalization and is not a purely syntactic process. By changing are some advantages to the current PUPS formulation. The formumantics of PUPS structures, the mechanism also has a well-workedlation of the mechanism is more uniform. Developed from the serules produced by the older ACT generalization mechanisms. There rules of broader generality. The production rules produced as comzation. It is a way of going beyond a particular single experience to pilations of the analogy process are, in fact, basically the generalized In the terminology of ACT*, analogy is a mechanism for general-

> eralizations - again, something subjects seem to have access to. able the subject to have conscious access to the basis for these gencontext. Also, by anchoring analogy in declarative structures, we en-ACT where the same generalization would emerge independent of

a generalized rule that was overapplying. required to look back and try to discover critical features to restrict examples work and so produce overgeneralizations. A process was the ACT* theory. Generalizations can miss critical features about why It was necessary to have a countervailing discrimination process in

amples are not adequately encoded in PUPS. There may be precona list of numbers: of the LISP call in FranzLISP (mapcar 'sub1 lis), which subtracts 1 from ditions to the successful operation of a rule that were missed in the is a little different. Overgeneralizations arise because the original exencoding. For instance, suppose a subject has the following encoding The same dilemma exists in the PUPS theory although its character

Example:

isa: function-call

function: (apply neg-op lis) form: (list mapcar' sub1 lis)

context: FranzLISP

isa: LISP.function

function: (implement neg-op) form: (text sub1)

but is working in INTERLISP. Below is the PUPS encoding of the new goal: Now suppose the person wants to add 1 to each number in a list

isa: function-call

function: (apply pos-op lis2) form: ???

context: INTERLISP

isa: LISP.function

function: (implement pos-op)

form: (text add1)

the example: The analogy process would calculate the following implication from

(list mapcar '(implement x) y) \rightarrow (apply x y),

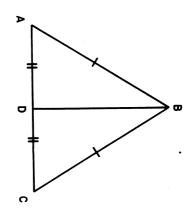
example. Therefore, the hypothesis is made that this is the critical addI, but these are already in the implication that led to the analogy some feature that cannot be mapped. There are features such as that FranzLISP and to the current goal that the context be INTERLISP feature, and we add as a precondition to example that the context be is FranzLISP in the former example and INTERLISP in the current the first example involved sub1 and the current example involved assume that the student does this by the same correspondence process cumstance where the student must try to make a discrimination. We The most immediate feature that does not map is the context, which in the example and the problem into correspondence, looking for that underlies making analogies. The student tries to put the elements the opposite order after the first attempt fails. Now we have a cirwe will suppose the student determines by experimenting and testing The implication for the example now becomes: Unfortunately, the correct argument order is (mapcar lis2 'add1) - as which leads to the inference that (mapcar 'add1 lis2) will do the trick

context (FranzLISP) & (list mapcar '(implement x) y) \rightarrow (apply x y)

into an INTERLISP-specific rule. slot) is available for future analogy, and that analogy can be compiled ogy is now blocked from going through in the future. On the other Given that INTERLISP cannot be mapped onto FranzLISP, the analhand, the INTERLISP example (with (list mapcar lis2 'add1) as a torm

the triangles were congruent. of the segment CD. Then he applied the reflexive rule to infer that student was finally able to apply the side-side-side rule to infer that reflexive to infer BD was congruent with itself. With this in place the that led nowhere. His final step in this floundering was to apply the AD was congruent to itself. These last two gave him legal inferences to infer that the measure of segment AD is congruent to the measure rules he knew would work. He applied the definition of congruence apply either directly. Then the student tried to apply some of the postulate and then the side-angle-side postulate but not to be able to atypical, was observed to think about using first the side-side-side side-side postulate and the side-angle-side postulate. One student, not tutor. Students when they encounter this problem know only the side-9.2 of a geometry problem that a student faced with the geometry ristic constraints to overgeneral rules. Consider the example in Figure A frequent context for performing discriminations is to add heu-

the reflexive rule – one involving AD, which had been unsuccess^{tul}, The student had created two examples for himself of the use of



Given: AD AB BC BC BC A

∆ABD ≅ ∆CBD

Figure 9.2. A geometry problem that invokes discrimination of the appropriate situation in which to employ the reflexive rule of congruence.

and one involving \overline{BD} , which had been successful. Let us consider crimination between the two of them: how the student represented these two inferences and formed a dis-

inference l

isa: geometry-inference

function: (help-prove goal)

form: (rule diagram reflexive statement1)

isa: geometry-statement

tunction: (to-be-proven problem1)

form: (statement triangle ABD congruent triangle CBD)

statement l

isa: geometry-statement

function: (conclusion-of inference1)

form: (statement segmentAD congruent segmentAD)

segmentAD

isa: segment

function: (part-of triangleABD)

form: (segment A D)

inference2 ısa: geometry-inference

function: (help-prove goal)
form: (rule diagram reflexive statement2)

statementz

isa: geometry statement function: (conclusion-of infere

function: (conclusion-of inference2) form: (statement segmentBD congruent segmentBD)

segmentBD

isa: segment

function: (shared-part triangle ABD triangle CBD)

form: (segment B D)

In comparing the successful inference2 with the unsuccessful inference1 the student would add the precondition that segment BD is shared by two triangles. This would then lead to a rule that would have the student infer that shared segments of triangles are congruent, and the application of this rule would not depend on the context in which the rule was evoked – side-side-side, side-angle-side, angle-angle-side, hypotenuse-leg, and so forth. This is in fact the behavior we observe of our students; their use of this rule is not restricted to proving triangles congruent by side-side-side as in this example.

The astute reader will note another feature of this rule, which is that it does not require (but should) the shared sides to be parts of to-be-proven congruent triangles. About half of the students we have observed appear to have induced the rule in the above form and will inappropriately apply the reflexive rule in one of the later problems that involved triangles with shared sides where the triangles are not to be proven congruent. From this, subjects learn an additional precondition by the same discrimination mechanism.

The interesting question concerns the other half of the students, who appear to learn this "part of to-be-proven congruent triangles" precondition from the first example, which does not seem to create the opportunity for learning it. That is, the successful and unsuccessful examples are not discriminated by this feature. We can only speculate, but it seems plausible to us that a good student might try to create an example that satisfies the precondition (shared parts of triangles) but does not serve the function (help to prove the problem). Such examples are easy to come by, and these self-generated examples would serve to force the desired discrimination.

It should be noted that even the constraint of shared segments of to-be-proven congruent triangles is not logically sufficient to guarantee that the rule is a necessary part of the proof. However, there are no proof problems that occur in high school geometry texts that

create situations that bring this out. So we have not been able to observe what high school students do. However, informally passing such problems among ourselves we have found that we have an irresistible urge at least covertly to make the reflexive inference. This seems to indicate that the acquisition of preconditions is not based on a formal analysis of logical necessity or sufficiency. Rather, as it is implemented in the PUPS system, it is based on empirical comparisons of cases where an example does or does not work.

A major feature of this discrimination mechanism is that the preconditions are stored as declarative embellishments along with the example. This means that a student can in principle examine them example. This potential for conscious filtering of proposed premadequate. This potential for conscious filtering of proposed preconditions makes the PUPS discrimination mechanism much more satisfactory than the ACT discrimination mechanism or other mechanisms, which are automatic comparison procedures whose results cannot be inspected because they immediately become embodied as productions.

Comparison to related work

It may be instructive to look at how our theory compares with other work on analogy and related topics. There are several similar projects.

Mitchell's generalization

cept (the goal concept), a single positive instance of the concept (the traincontext. The system is given a high-level description of the target cona concept based on a single example, and without a problem-solving Mitchell (Mitchell, Keller, & Kedar-Cabelli, 1986) forms descriptions of ing example), a description of what an acceptable concept definition simply expands the terms in a high-level description until all the terms An EBG (explanation-based generalization) algorithm tries to find a would be (the operationality criterion), and a list of facts about the domain. training example satisfies the goal concept, the proof is generalized to sion of graspable will be has-handle). After a proof is generated that the small-diameter, and the training example is has-handle, then the expanexample is used (for example, if graspable means either has-handle or in the description meet the operationality criterion. If disjuncts are ever proof that the training example satisfies the goal concept. To do this it Included in these facts are abstract rules of inference about the domain. found in a term's expansion, the disjunct corresponding to the training

form a rule that is capable of matching any instance of the goal concept that meets this same low-level description. Note that this generalization is in general more restrictive than the goal concept, reflecting the choices made at any disjuncts encountered during the expansion process. Since the entire tree of expansions is saved during this process, a side effect is that the system can explain why the training example is an instance of the goal concept: It can point to specific features of the example that fulfill the various criteria specified in the high-level concept description (Mitchell et al., 1986).

The expansions done by the EBG method are not unlike the elaborations done by the PUPS system. The essential difference is that, whereas the EBG system blindly expands until it reaches a dead end or the (apparently ad hoc) operationality criterion is met, the PUPS system has an implicit operationality criterion, which is that the expansion is sufficiently elaborate for the no-function-in-content principle to apply. A second difference is that PUPS need not be given abstract rules of inference for the domain. It tries to infer these directly from its encoding of examples. Thus EBG starts out with a strong domain theory and essentially composes new rules; whereas PUPS discovers the rules hidden in its examples. A third difference is that the EBG method simply characterizes the way in which a single object instantiates a concept, whereas PUPS draws analogies in order to further problem-solving efforts.

Kedar-Cabelli's purpose-directed analogy

Kedar-Cabelli (1985) developed an analogy system based on EBG. Her system is typically given an object description and the task of answering a question such as "Does this object serve function x?" In order to answer the question, the system searches its knowledge base for model objects that serve function x and tries to "prove" that the model object serves the function. Then it tries to use this proof to prove that the target object serves the specified function. If the proof is successful, a generalization is formed that could essentially answer this question directly, without the analogy process.

Kedar-Cabelli's system is like ours with the differences already noted with respect to Mitchell's system. It is basically the obvious application of EBG to analogy. The major difference in the way it treats analogy is that whereas she asks questions of the form "Is it true that A serves function x?" we ask questions of the form "What function does A serve?" Another difference is the extent of the elaboration; Kedar-Cabelli elaborates her examples until she cannot elaborates.

orate further, whereas PUPS simply elaborates as far as is necessary for the no-function-in-content principle to apply.

Winston's analogy

example, making inferences about physical properties that are not single positive instance. The system elaborates the description of the off with a description of the target concept and a description of a additional function of providing rules of inference. The system starts examples and not stored separately, and thus the example serves the situation similar to the current one. of the concept, is built in the form of a rule that could apply in a scribing the way in which the current example fulfills the description is stable as well and that the purpose of its flat bottom is stability various features of the example. For example, if the target example this using "precedents" in order to provide functional descriptions of example meets the description of the target concept. ANALOGY does the system attempts to show that the elaborated description of the the elaboration is complete (i.e., no more elaboration can take place), explicitly represented in the input description of the example. Once major difference is that the rules of inference are stored with the Assuming that this proof attempt is successful, a generalization, dethey have flat bottoms, then the system concludes that the example has a flat bottom and the system knows that bricks are stable because 1983) is very similar to the work of Mitchell and Kedar-Cabelli. The Winston's ANALOGY system (Winston, Binford, Katz, & Lowry,

One of the big differences between ANALOGY and PUPS is that ANALOGY seems capable only of filling in function slots. That is, ANALOGY would not be able to generate an example that served a specific function. Indeed, this observation could be made of all the research we have looked at, except for Carbonell's work (which is capable only of finding form that fills a particular function and not the reverse). PUPS is the only system we know of that can draw analogies in either direction.

Rumelhart and Norman's analogy

Rumelhart and Norman (1981) discuss a method for drawing analogies that consists of generating a description of the differences between the model and the goal and writing new form that observes those differences. Their algorithm would compare functional descriptions of the model and the goal and notice that the goal is "just like"

for making the appropriate substitution of y wherever an x appeared algorithm would then copy the structure of the model exactly, except the target, except that you use x for y." Using this information, the in the model.

it is difficult to know exactly what the termination condition of the eralization process. They also do not discuss elaboration in detail, so of model selection, and they make no mention of any kind of genelaboration process is melhart and Norman's model. There is no discussion of the problems are a number of things in PUPS that correspond to nothing in Ru-What is specified in this model is very like our own work, but there

Gentner's structure-mapping

description of the domain are much more likely to be mapped than aticity principle says that those relations that are central to the functional sun (e.g., hot, yellow) do not. An analogy that maps a large number system to an atom. For instance, the relationship between electrons of two or more arguments. In an analogy, one is concerned only with anything). This causal analysis is similar to Winston's (1979) model. assertion that the sun is hotter than the planets (which doesn't cause relation is more likely to be mapped to the domain of atoms than the than a planet in some way causes the planet to orbit the sun. Thus this those that are not. So, for instance, the fact that the sun is more massive target domain involves a causal analysis of the domains. The systemabstraction. The method for selection of what features will map to the of attributes is called a literal similarity. An analogy in which the model and the nucleus should be mapped, but the features specific to the ural way those features that should map when comparing the solar mapping relations. From this assumption, she distinguishes in a natare predicates taking one argument, and relations, which are predicates types of features of the model. In particular, there are attributes, which is an abstract description rather than a physical object is called an Gentner's (1983, this volume) theory distinguishes among various be mapped, then you should not map it. The central idea is that, if you cannot show a reason for a relation to

Carbonell's derivational analogy

discussed. His basic strategy is to take a worked-out solution for a Carbonell's (1985) work is different in kind than the systems so far problem and convert it to the current task. The problem solution may

> are then evaluated with respect to the current problem, and various operators and the parts of the problem domain. These dependencies along with an elaborate description of the dependencies among the exist at one level, they are combined by a learning/generalization leting old ones. When enough related instances of a problem solution changing the order of the operators, inserting new operators, or depropriate for the current problem. The editing operations include editing operations are performed to convert the solution to one appoints along the problem-solving continuum) as a list of operators be represented at any level of abstraction (corresponding to various process into a more general solution.

a solution to the current problem. In our work, each operator applisolution be transported (modulo certain possible transformations) into or an analogy), and our solutions may therefore potentially borrow cation is done by a separate step (which may be either a learned rule he represents problem solutions as a whole and requires that the entire more piecemeal approach is closer to the human use of analogy. (our theory predicts more transfer to novel problems). We think this the entire solution, these generalizations are more widely applicable learn describe an individual step in a problem solution rather than from many different examples. Also, since the generalizations we A major difference between Carbonell's work and our own is that

Conclusion

processing: A general framework for analogy might have the following steps of

- Obtain a goal problem.
- Find an example similar to the problem
- Elaborate the goal.
- Generate a mapping between the goal and the example
- Use the mapping to fill in the goal pattern. Check the validity of the solution.
- Generalize and form a summarization rule

with an elaboration of the example. The consequence of this is that accomplish the steps and in the order in which the steps are done this framework. The differences between the systems lie in how they and, By contrast, Winston's (Winston et. al., 1983) system does the For instance, in PUPS the elaboration of the goal is done in parallel It is apparent that the systems we have discussed by and large ht with the mapping is a by-product of the elaboration and thus trivial to

mapping explicitly. elaboration before it searches for the example and must generate the

serves that function (in terms of a problem solution). PUPS, by conof a goal that needs to be accomplished) and provides a structure that trast, can do both of these tasks. function. Carbonell's system is given a functional description (in terms example form and determine whether or not it serves a particular (such as Winston's and Mitchell's) in which the object is to take an the kind of questions the systems attempt to answer. There are systems Another way in which we can contrast the systems is by looking at

NOTES

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- 2 This discussion differs slightly from the implementation.

 The logical justification for this derives from the fact that the relationship among multiple functions is assumed to be biconditional - if one function is satisfied, then all are.

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